

ZGC Concurrent Class Unloading

Another safepoint operation bites the dust

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Erik Österlund Garbage Collection Engineer Java Platform Group, Oracle February 04, 2019

Live for the Coce



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Agenda

- 1 What is ZGC? What is class unloading?
- 2 Overview of phases
- 3 Concurrent code and metadata unloading
- 4 Evaluation
- **5** Future plans



What is ZGC?



New Concurrent GC in JDK 11

(Experimental feature, Linux/x86_64 only)



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ZGC Goals

Multi-terabyte heaps





TB

Easy to tune



15% Max application throughput reduction



What is a concurrent GC

	Serial	Parallel	CMS	G1	ZGC
Marking	-	-	*	*	
Relocation/Compaction	-	-	-	-	
Reference Processing	-	-	-	-	
Relocation Set Selection	-	-	-	-	
StringTable Cleaning	-	-	-	-	
JNI WeakRef Cleaning	-	-	-	-	
JNI GlobalRefs Scanning	-	-	-	-	
Class Unloading	-	-	-	-	
Thread Stack Scanning	-	-	-	-	-



Concurrent Class Unloading Released in JDK 12 for ZGC



Traditional Class Unloading

- Step 1: Marking (concurrent)
 - Mark metadata (classes, CLDs) when marking objects
- Step 2: Reference processing (STW)
 - Need to know what is reachable from finalizers before class unloading
- Step 3: Unloading (STW)
 - Unload code cache
 - Unload metadata



ZGC Concurrnet Class Unloading

- Step 1: Marking (concurrent)
 - Mark metadata (classes, CLDs) when marking objects
 - Mark both strong and final reachable graphs
- Step 2: Reference processing (concurrent)
 - Already know what is reachable from finalizers before class unloading
- Step 3: Unloading (concurrent)
 - Unload code cache
 - Unload metadata









Step 1: Marking



Marking overview

- Color heap object pointers with appropriate marked color
 - Special bit pattern for edges reachable from finalizers only
 - Mutator load barriers upgrade them to strongly reachable when loaded
- Mark metadata objects similarly
 - Mark metadata reachable from objects, with strong/final reachability



Metadata graph





Step 2: Reference Processing



Reference processing overview

- WeakReferences cleared if referent not strongly reachable
- PhantomReferences cleared if referent not reachable
- "Weak" VM datastructures have "phantom" strength — Classes die if not reachable (including from finalizers)
- Each access on weak/phantom is annotated using my Access API
- A class is dead if a phantom load of its holder returns NULL



Step 3: Unloading



Stale Datastructures after Reference Processing

Subklass/sibling/implementor lists Method data objects Instance class dependency context jli.CallSite dependency context Class loader data graph Protection domain cache table Module table String table Package table Symbol table Resolved method table Loader constraint table Resolution error table Metaspace Code cache (JIT compiled code) Inline caches





Basically everything is a huge mess... ...and we just continue running anyway



ZGC Unloading Overview

Pause Mark End





ZGC Unloading Overview

Pause Mark End





ZGC Unloading Overview

Pause Mark End





Concurrent Code Unloading



Code cache

- Colored pointers into Java heap
 - Misaligned immediate values
 - Which color should native compiled object references have?
 - Need a way to paint native compiled object references
- Inline caches (CompiledIC) pointing at now dead native methods (nmethods), because of dead object references
 - Running any such code yields crashes
 - Need a way of preventing calls to dead native methods





NMethod Entry Barriers

- Arm all nmethods not on stack in GC pause
 - Change global epoch value, caught with cmp; je; at verified entry
- Trap calls to armed nmethods
 - NMethods are "good" or "bad" based on object pointer liveness
- When entering good nmethods
 - Fix up object pointers (oops)
 - Disarm barrier by patching cmp immediate value
- When entering bad nmethods
 - Re-resolve the call





Example: Calling inline cache to dead nmethod











Example: Calling inline cache to stale but live nmethod











Concurrent Code Unloading

• Unlink

- Clean stale inline caches (patch machine code that Java threads run)
- Fixup object references (patching more machine code)
- Disarm entry barriers (yet some more machine code patching)
- Unlink nmethods from dependency contexts (lock-free unlinking)
- Unlink exception caches (more lock-free unlinking)
- Global rendezvous handshake
- Purge
 - Purge unloading nmethods with make_unloaded()
 - Sweeper subsequently frees up memory in code cache



Concurrent Metadata Unloading



Structure

- Unlink
 - Expose logically already unlinked view of data to mutators
 - Subclass/sibling/implementor lists (lock-free)
 - Method data object (lock-free and per-MDO lock)
 - Protection domain cache (lock)
 - Class loader data graph (lock)
 - StringTable and SymbolTable (crazy concurrent)
- Rendezvous handshake
- Purge
 - Delete Klass, Method, CLD, handles, table entries, etc.



Evaluation



SPECjbb[®]2015 – Score



Mode: Composite Heap Size: 128G OS: Oracle Linux 7.4 HW: Intel Xeon E5-2690 2.9GHz 2 sockets, 16 cores (32 hw-threads)

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SPECjbb[®]2015 – Pause Times





SPECjbb[®]2015 – Pause Times





SPECjbb[®]2015 – Pause Times









- Short-term
 - Turn ZGC into a product feature





- Short-term
 - Turn ZGC into a product feature

Remove experimental status





- Short-term
 - Turn ZGC into a product feature





- Short-term
 - Turn ZGC into a product feature

Long-term

- Generational
- Sub-millisecond max pause times
- Additional platform support
- Graal JIT support





- Short-term
 - Turn ZGC into a product feature

Generational

- Withstand higher allocation rates
- Lower heap overhead
- Lower CPU usage

• Long-term

- Generational
- Sub-millisecond max pause times
- Additional platform support
- Graal JIT support





- Short-term
 - Turn ZGC into a product feature

• Long-term

- Generational
- Sub-millisecond max pause times
- Additional platform support
- Graal JIT support

Sub-millisecond max pause times

- Within reach
- Reduce root set size
- Time-to-Safepoint, etc





- Short-term
 - Turn ZGC into a product feature

Long-term

- Generational
- Sub-millisecond max pause times
- Additional platform support
- Graal JIT support

Additional platform support

- macOS?
- Windows?
- Sparc?
- Aarch64?





- Short-term
 - Turn ZGC into a product feature

Long-term

- Generational
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Get Involved!



ZGC Project

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zgc-dev@openjdk.java.net

http://wiki.openjdk.java.net/display/zgc/Main



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ZGC Project

Source Code





http://hg.openjdk.java.net/jdk/jdk

http://hg.openjdk.java.net/zgc/zgc



Thanks!



Questions?



